## On the Period Length of Pseudorandom Vector Sequences Generated by Matrix Generators

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Abstract. In Tahmi [5], Niederreiter [4], Afflerbach and Grothe [1], and Grothe [2] linear recursive congruential matrix generators for generating *r*-dimensional pseudorandom vectors are analyzed. In particular, conditions are established which ensure that the period length equals  $p^r - 1$  for any nonzero starting vector in case of a prime modulus p. For a modulus of the form  $p^{\alpha}$ ,  $\alpha \geq 2$  and p prime, this paper describes a simple method for constructing matrix generators having the maximal possible period length  $(p^r - 1) \cdot p^{\alpha - 1}$  for any starting vector which is nonzero modulo p.

1. Introduction and Notation. A linear recursive congruential matrix generator for generating r-dimensional pseudorandom vectors is of the form

(1) 
$$\vec{x}_{n+1} \equiv A \cdot \vec{x}_n \pmod{m}, \quad \vec{x}_{n+1} \in \mathbf{Z}_m^r, \ n \ge 0,$$

where the modulus m is a positive integer,  $\mathbf{Z}_m = \{0, 1, \ldots, m-1\}, \ \vec{x}_0 \in \mathbf{Z}_m^r$ , and  $A \in \mathbf{Z}_m^{r \times r}$ , i.e., A is an  $r \times r$ -matrix with elements in  $\mathbf{Z}_m$ . In the sequel it is assumed that the matrix A is nonsingular modulo m. Then the vector sequence  $(\vec{x}_n)_{n\geq 0}$  generated by (1) is purely periodic, and the smallest positive integer  $\lambda = \lambda(A, \vec{x}_0, m)$  with  $\vec{x}_{\lambda} = \vec{x}_0$  is called the *period length of the vector sequence*  $(\vec{x}_n)_{n\geq 0}$ . Analogously, the matrix sequence  $(A_n)_{n\geq 0}$  with  $A_n \equiv A^n \pmod{m}, A_n \in \mathbf{Z}_m^{r \times r}$ , is purely periodic, and the smallest positive integer  $\lambda = \lambda(A, m)$  for which  $A_{\lambda}$  equals the identity matrix I is called the *period length of the matrix sequence*  $(A_n)_{n\geq 0}$ . The following two remarks are immediate consequences of these definitions.

Remark 1. The period length  $\lambda(A, \vec{x}_0, m)$  of the vector sequence  $(\vec{x}_n)_{n\geq 0}$  divides the period length  $\lambda(A, m)$  of the matrix sequence  $(A_n)_{n\geq 0}$  for any starting vector  $\vec{x}_0 \in \mathbf{Z}_m^r$ .

Remark 2. If  $A_{\nu} = I$  for some positive integer  $\nu$ , then the period length  $\lambda(A, m)$  of the matrix sequence  $(A_n)_{n\geq 0}$  divides  $\nu$ .

It is well known (cf. Tahmi [5], Niederreiter [4], and Grothe [2]) that  $\lambda(A, \vec{x}_0, p) = \lambda(A, p) = p^r - 1$  for any starting vector  $\vec{x}_0 \in \mathbf{Z}_p^r \setminus \{\vec{0}\}$  in case of a prime modulus m = p if the characteristic polynomial of the matrix A is primitive modulo p. In this paper the case of a modulus  $m = p^{\alpha}$ ,  $\alpha \ge 2$ , is considered where p is a prime number. It is shown that for  $p \ge 3$  or  $r \ge 2$  there exist matrix generators (1) with period length  $(p^r - 1) \cdot p^{\alpha - 1}$  for any starting vector which is nonzero modulo p, and a simple method is described for determining such a generator. Observe that  $(p^r - 1) \cdot p^{\alpha - 1}$  is the maximal possible period length according to the following technical lemma.

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## 2. Matrix Generators with Maximal Period Length.

LEMMA. Let  $A \in \mathbb{Z}_{p^{\alpha+1}}^{r \times r}$ ,  $\alpha \geq 1$ , be a matrix which is nonsingular modulo p, and define matrices  $A_n \in \mathbb{Z}_{p^{\alpha+1}}^{r \times r}$  by  $A_n \equiv A^n \pmod{p^{\alpha+1}}$ ,  $n \geq 0$ . Let  $\lambda_{\alpha} = \lambda(A, p^{\alpha})$  and  $\lambda_{\alpha+1} = \lambda(A, p^{\alpha+1})$  denote the period lengths of the matrix sequence  $(A_n)_{n\geq 0}$  modulo  $p^{\alpha}$  and modulo  $p^{\alpha+1}$ , respectively. Then

$$\lambda_{\alpha+1} = \begin{cases} \lambda_{\alpha} & \text{for } A^{\lambda_{\alpha}} \equiv I \pmod{p^{\alpha+1}}, \\ \lambda_{\alpha} \cdot p & \text{for } A^{\lambda_{\alpha}} \neq I \pmod{p^{\alpha+1}}. \end{cases}$$

*Proof.* From  $A^{\lambda_{\alpha}} \equiv I \pmod{p^{\alpha}}$  it follows that  $A^{\lambda_{\alpha}} = I + p^{\alpha} \cdot B$  for some matrix  $B \in \mathbb{Z}^{r \times r}$ . Therefore,

$$A^{\lambda_{\alpha}\cdot p} = (I + p^{\alpha} \cdot B)^{p} = I + {p \choose 1} \cdot p^{\alpha} \cdot B + {p \choose 2} \cdot (p^{\alpha} \cdot B)^{2} + \dots + (p^{\alpha} \cdot B)^{p},$$

which yields  $A^{\lambda_{\alpha} \cdot p} \equiv I \pmod{p^{\alpha+1}}$ , i.e.,  $\lambda_{\alpha+1}$  divides  $\lambda_{\alpha} \cdot p$  according to Remark 2. From  $A^{\lambda_{\alpha+1}} \equiv I \pmod{p^{\alpha+1}}$  it follows that  $A^{\lambda_{\alpha+1}} \equiv I \pmod{p^{\alpha}}$ , i.e.,  $\lambda_{\alpha}$  divides  $\lambda_{\alpha+1}$  according to Remark 2, which proves the lemma.  $\Box$ 

The purpose of this paper is to prove the following result.

THEOREM. Let  $B \in \mathbb{Z}_{p^{\alpha}}^{r \times r}$ ,  $\alpha \geq 2$ , be a matrix whose characteristic polynomial is primitive modulo p. Then

(2) 
$$B^{p^r-1} \equiv I + p \cdot C \pmod{p^2}$$

for some matrix  $C \in \mathbf{Z}_p^{r \times r}$ . Let  $D \in \mathbf{Z}_{p^{\alpha-1}}^{r \times r}$  denote an arbitrary matrix with  $B \cdot D \equiv D \cdot B \pmod{p}$ ,

(3) 
$$\det(D) \not\equiv 0 \pmod{p} \quad for \ p \ge 3,$$

and

(4) 
$$\det(D) \equiv \det(D+I) \equiv 1 \pmod{2} \quad for \ p=2.$$

Define a matrix  $A \in \mathbf{Z}_{p^{\alpha}}^{r \times r}$  by

,

(5) 
$$A \equiv B \cdot (I + p \cdot (C - D)) \pmod{p^{\alpha}}$$

Then the period length of the vector sequence  $(\vec{x}_n)_{n\geq 0}$  generated according to (1) with matrix A and modulus  $m = p^{\alpha}$  is given by

 $\lambda(A, \vec{x}_0, p^{\alpha}) = (p^r - 1) \cdot p^{\alpha - 1}$ 

for any starting vector  $\vec{x}_0 \in \mathbf{Z}_{p^{\alpha}}^r$  with  $\vec{x}_0 \not\equiv \vec{0} \pmod{p}$ .

*Proof.* The proof is subdivided into four parts (i) to (iv).

(i) Because of  $A \equiv B \pmod{p}$  according to (5) it follows that

(6) 
$$\lambda(A, \vec{x}_0, p) = \lambda(A, p) = p^r - 1$$

for any starting vector  $\vec{x}_0 \in \mathbb{Z}_p^r \setminus \{\vec{0}\}$ , since the characteristic polynomial of the matrix B is primitive modulo p. In particular,  $B^{p^r-1} \equiv I \pmod{p}$  holds. Hence a matrix C with (2) exists. Observe that (2) yields  $B \cdot C \equiv C \cdot B \pmod{p}$ , which implies that  $B \cdot (C-D) \equiv (C-D) \cdot B \pmod{p}$  because of the hypothesis  $B \cdot D \equiv D \cdot B \pmod{p}$ . Therefore (5) and (2) yield

(7) 
$$A^{p^{r}-1} \equiv [B \cdot (I+p \cdot (C-D))]^{p^{r}-1} \equiv B^{p^{r}-1} \cdot (I+(p^{r}-1) \cdot p \cdot (C-D))$$
$$\equiv (I+p \cdot C) \cdot (I-p \cdot (C-D)) \equiv I+p \cdot D \pmod{p^{2}}.$$

If p = 2 then it follows from (7) that

$$A^{2^r-1} = I + 2 \cdot D + 4 \cdot E$$

for some matrix  $E \in \mathbf{Z}^{r \times r}$  and hence

$$A^{(2^{r}-1)\cdot 2} = (I+2\cdot D+4\cdot E)^{2} = I+4\cdot D+4\cdot D^{2}+8\cdot F$$

for some matrix  $F \in \mathbf{Z}^{r \times r}$ , i.e.,

$$A^{(2^r-1)\cdot 2} \equiv I + 4 \cdot D \cdot (D+I) \pmod{8}.$$

(ii) Now it is shown by induction that in case of  $p \ge 3$ ,

(8) 
$$A^{(p^r-1)\cdot p^{\nu}} \equiv I + p^{\nu+1} \cdot D \pmod{p^{\nu+2}}$$

for  $0 \le \nu \le \alpha - 2$ . Obviously, (7) is equivalent to (8) for  $\nu = 0$ . If (8) is valid for some  $\nu$  with  $0 \le \nu \le \alpha - 3$ , then

$$A^{(p^{r}-1)\cdot p^{\nu}} = I + p^{\nu+1} \cdot D + p^{\nu+2} \cdot E_{\nu}$$

for some matrix  $E_{\nu} \in \mathbf{Z}^{r \times r}$  and hence

$$A^{(p^{r}-1)\cdot p^{\nu+1}} = (I+p^{\nu+1}\cdot (D+p\cdot E_{\nu}))^{p} = I+p^{\nu+2}\cdot (D+p\cdot E_{\nu})+p^{\nu+3}\cdot F_{\nu}$$

for some matrix  $F_{\nu} \in \mathbf{Z}^{r \times r}$  because of  $p \geq 3$ , which yields

$$A^{(p^{r}-1)\cdot p^{\nu+1}} \equiv I + p^{\nu+2} \cdot D \pmod{p^{\nu+3}}.$$

Therefore (8) holds for  $0 \le \nu \le \alpha - 2$ . It can be similarly proved that in case of p = 2,

(9) 
$$A^{(2^{r}-1)\cdot 2^{\nu}} \equiv I + 2^{\nu+1} \cdot D \cdot (D+I) \pmod{2^{\nu+2}}$$

for  $1 \leq \nu \leq \alpha - 2$ .

(iii) Because of (3), (4), (6), (7), (8) and (9) it follows from the lemma that

(10) 
$$\lambda(A, p^{\nu+1}) = (p^r - 1) \cdot p^{\nu}$$

for  $0 \le \nu \le \alpha - 1$ . Note that if  $\vec{x}_0 \not\equiv \vec{0} \pmod{p}$ , then

$$D \cdot \vec{x}_0 \not\equiv \vec{0} \pmod{p} \quad \text{for } p \ge 3$$

and

$$D \cdot (D+I) \cdot \vec{x}_0 \not\equiv \vec{0} \pmod{p} \quad \text{for } p = 2$$

because of (3) and (4), respectively. Therefore (7), (8) and (9) show that

(11) 
$$A^{(p^r-1)\cdot p^{\nu}} \cdot \vec{x}_0 \not\equiv \vec{x}_0 \pmod{p^{\nu+2}}$$

for  $\vec{x}_0 \not\equiv \vec{0} \pmod{p}$  and  $0 \leq \nu \leq \alpha - 2$ .

(iv) Now it is proved by induction that

(12) 
$$\lambda(A, \vec{x}_0, p^{\nu+1}) = (p^r - 1) \cdot p^{\nu}$$

for any starting vector  $\vec{x}_0 \in \mathbf{Z}_{p^{\alpha}}^r$  with  $\vec{x}_0 \not\equiv \vec{0} \pmod{p}$  and  $0 \leq \nu \leq \alpha - 1$ . Obviously, (6) is equivalent to (12) for  $\nu = 0$ . Now assume that (12) is valid for some  $\nu$  with  $0 \leq \nu \leq \alpha - 2$ . Then

$$\lambda(A, \vec{x}_0, p^{\nu+2}) = \mu \cdot (p^r - 1) \cdot p^{\nu}$$

for some integer  $\mu \geq 1$ . Since

$$\lambda(A, \vec{x}_0, p^{\nu+2}) \neq (p^r - 1) \cdot p^{\nu}$$

according to (11), it follows that  $\mu > 1$ . Remark 1 and (10) imply that  $\lambda(A, \vec{x}_0, p^{\nu+2})$  divides  $(p^r - 1) \cdot p^{\nu+1}$  and hence  $\mu = p$ , which proves the theorem.  $\Box$ 

Observe that there exist primitive polynomials of degree r over the Galois field GF(p) for every positive integer r and every prime number p. Such a polynomial, and hence a matrix  $B \in \mathbb{Z}_{p^{\alpha}}^{r \times r}$  which satisfies the hypothesis of the theorem, can be determined without any effort if p and r are small integers (see, e.g., Knuth [3, p. 28]).

Since the characteristic polynomial of the matrix B is primitive modulo p, it follows that  $\det(B) \not\equiv 0 \pmod{p}$  and that  $B \cdot \vec{x}_0 \not\equiv \vec{x}_0 \pmod{2}$  for  $p = 2, r \ge 2$ , and  $\vec{x}_0 \not\equiv \vec{0} \pmod{2}$ . Hence  $\det(B+I) \equiv 1 \pmod{2}$  for p = 2 and  $r \ge 2$ . Therefore, the matrix  $D \in \mathbb{Z}_{p^{\alpha-1}}^{r \times r}$  with  $D \equiv B \pmod{p^{\alpha-1}}$  satisfies the hypothesis of the theorem if  $p \ge 3$  or  $r \ge 2$ .

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